



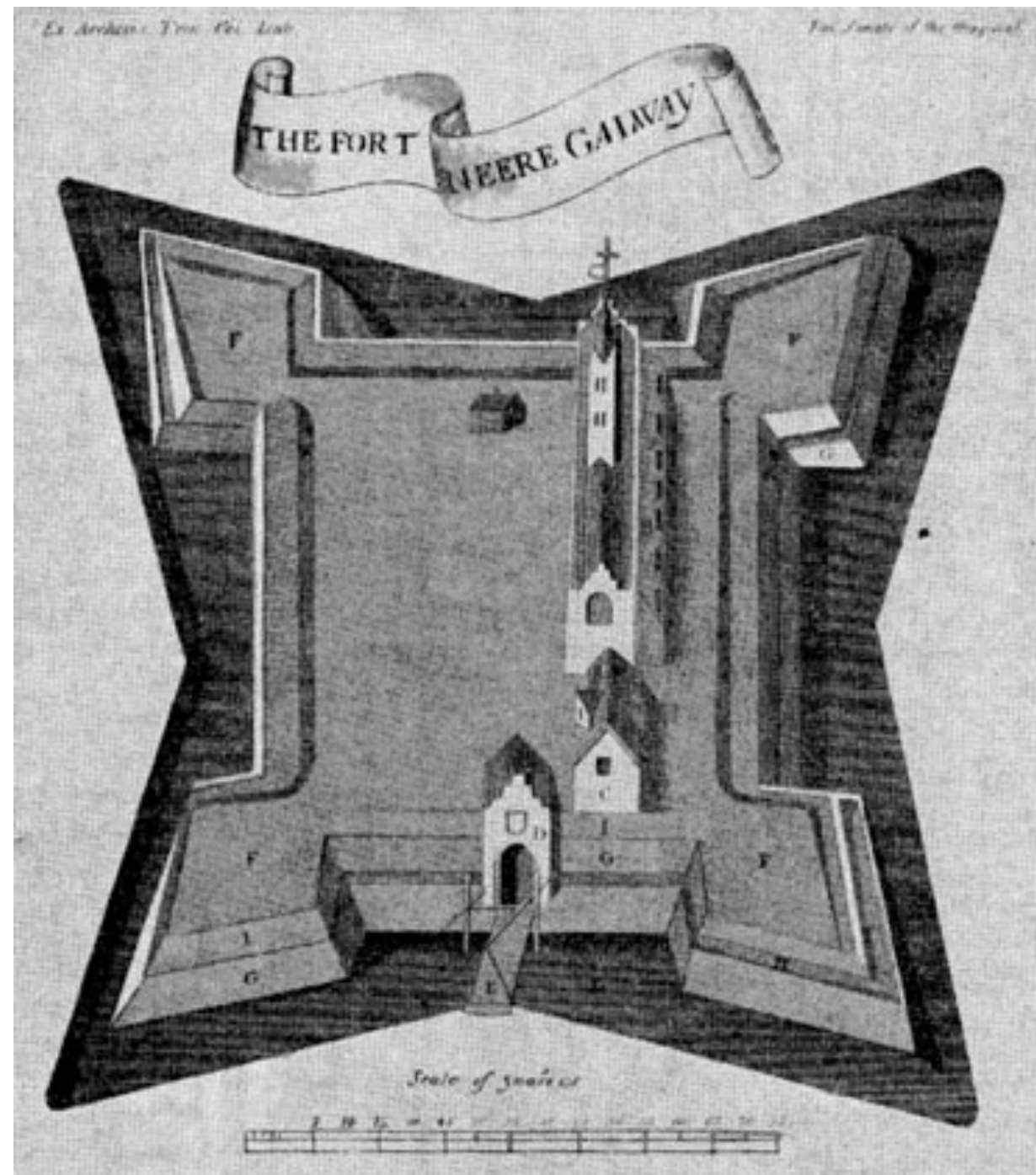
PennState

# CSE543: Computer Security

## Module: Web Security

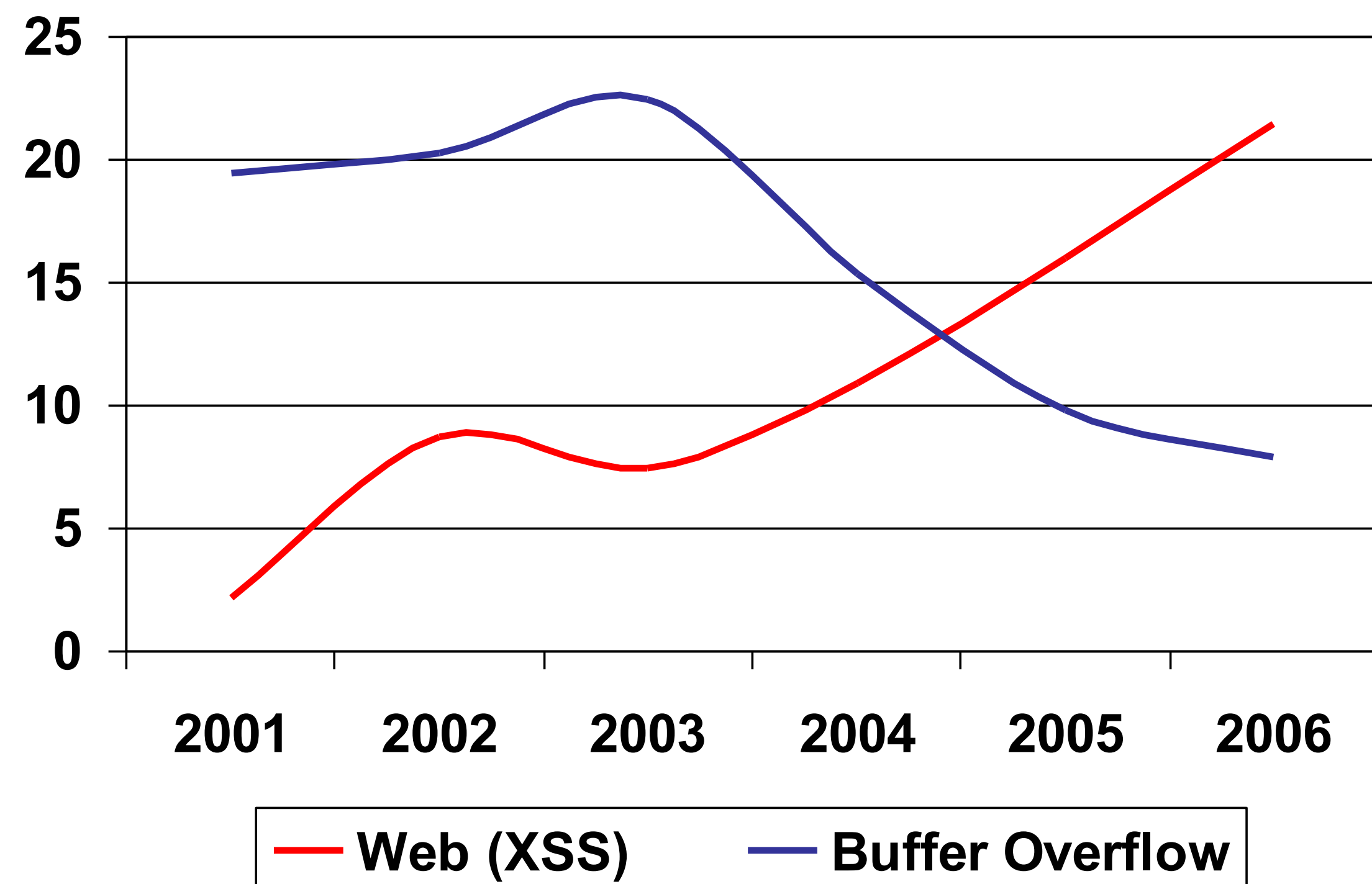
Prof. Syed Rafiul Hussain  
Department of Computer Science and Engineering  
Pennsylvania State University

# Network vs. Web Security



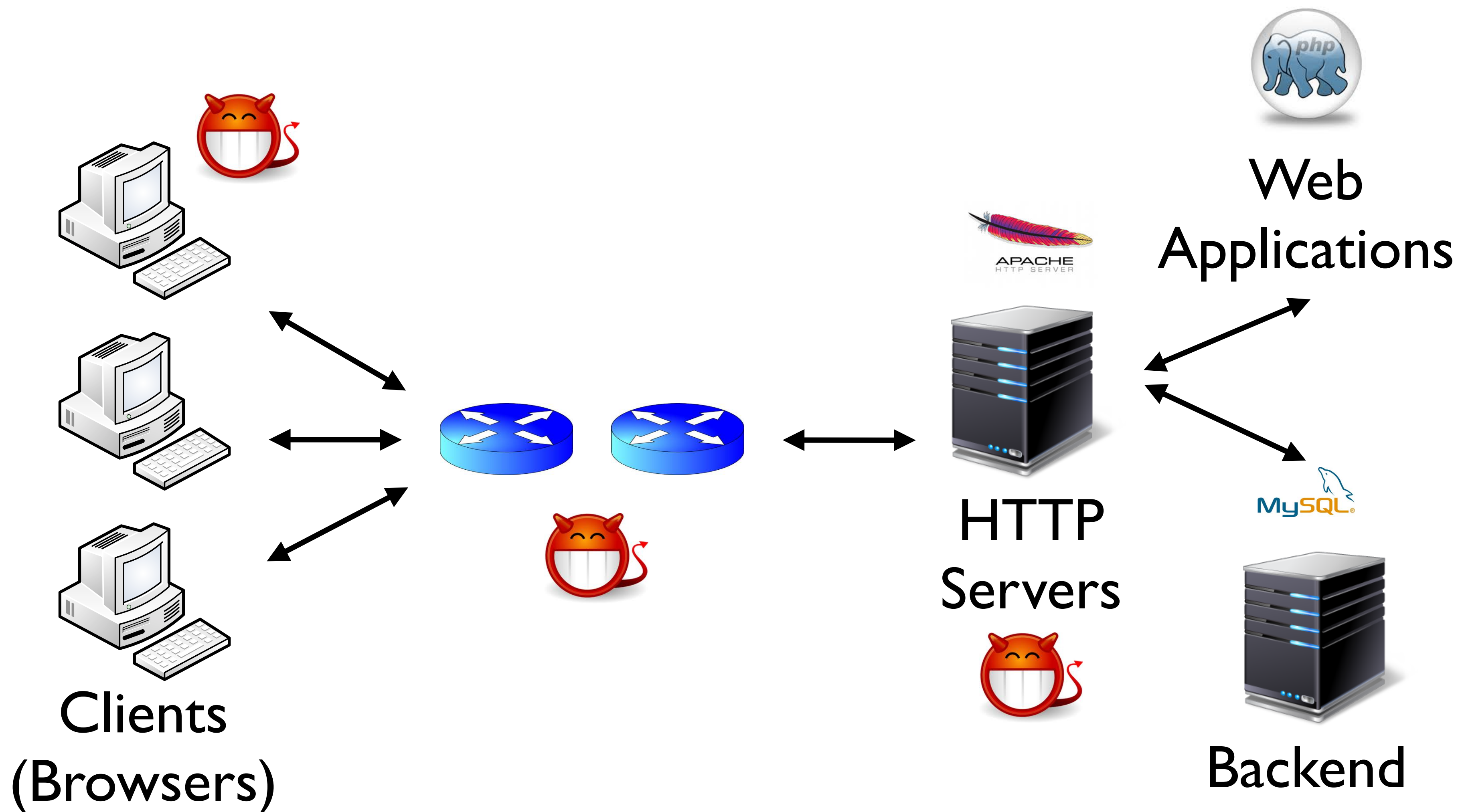
# Web Vulnerabilities

- Web vulnerabilities surpassed OS vulnerabilities around 2005
  - ▶ The “new” buffer overflow



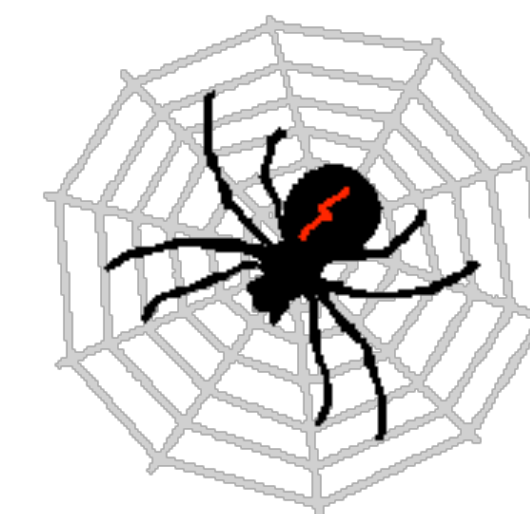
# Components of the Web

- Multiple interacting components



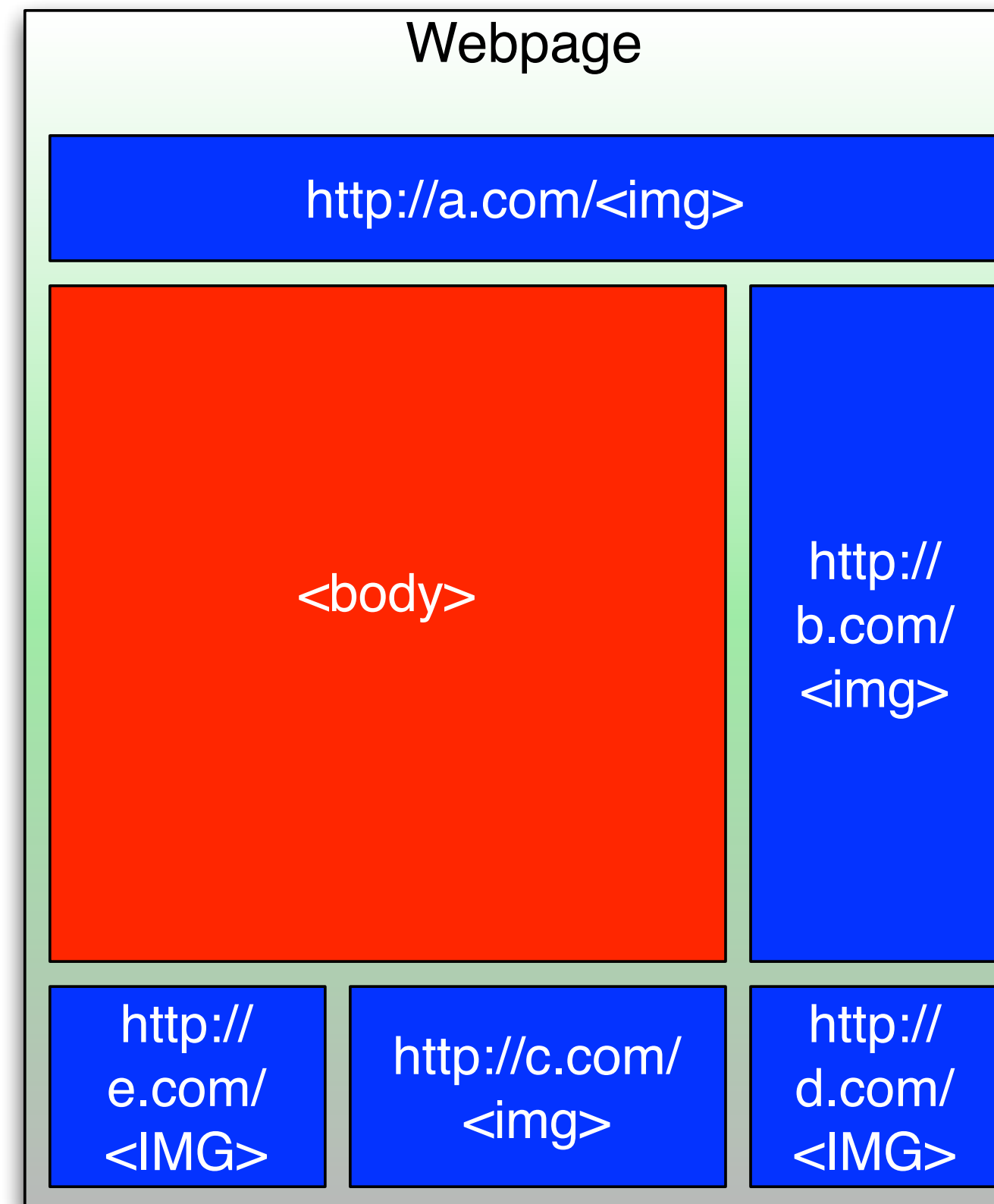
# Web security: the high bits

- The largest distributed system in existence
- Multiple sources of threats, varied threat models
  - ▶ Users
  - ▶ Servers
  - ▶ Web Applications
  - ▶ Network infrastructure
  - ▶ We shall examine various threat models, attacks, and defenses
- Another way of seeing web security is
  - ▶ Securing the web **infrastructure** such that the **integrity, confidentiality, and availability** of content and user information is maintained



# Early Web Systems

- Early web systems provided a click-render-click cycle of acquiring web content.
  - ▶ Web content consisted of static content with little user interaction.



- Browser sends HTTP requests to the server
  - ▶ Methods: GET, POST, HEAD, ...
  - ▶ GET: to retrieve a resource (html, image, script, css,...)
  - ▶ POST: to submit a form (login, register, ...)
  - ▶ HEAD (a HEAD request could its Content-Length header to check the filesize without actually downloading the file)
- Server replies with a HTTP response
- Stateless request/response protocol
  - ▶ Each request is independent of previous requests
  - ▶ Statelessness has a significant impact on design and implementation of applications

# Adding State to the Web: Cookies

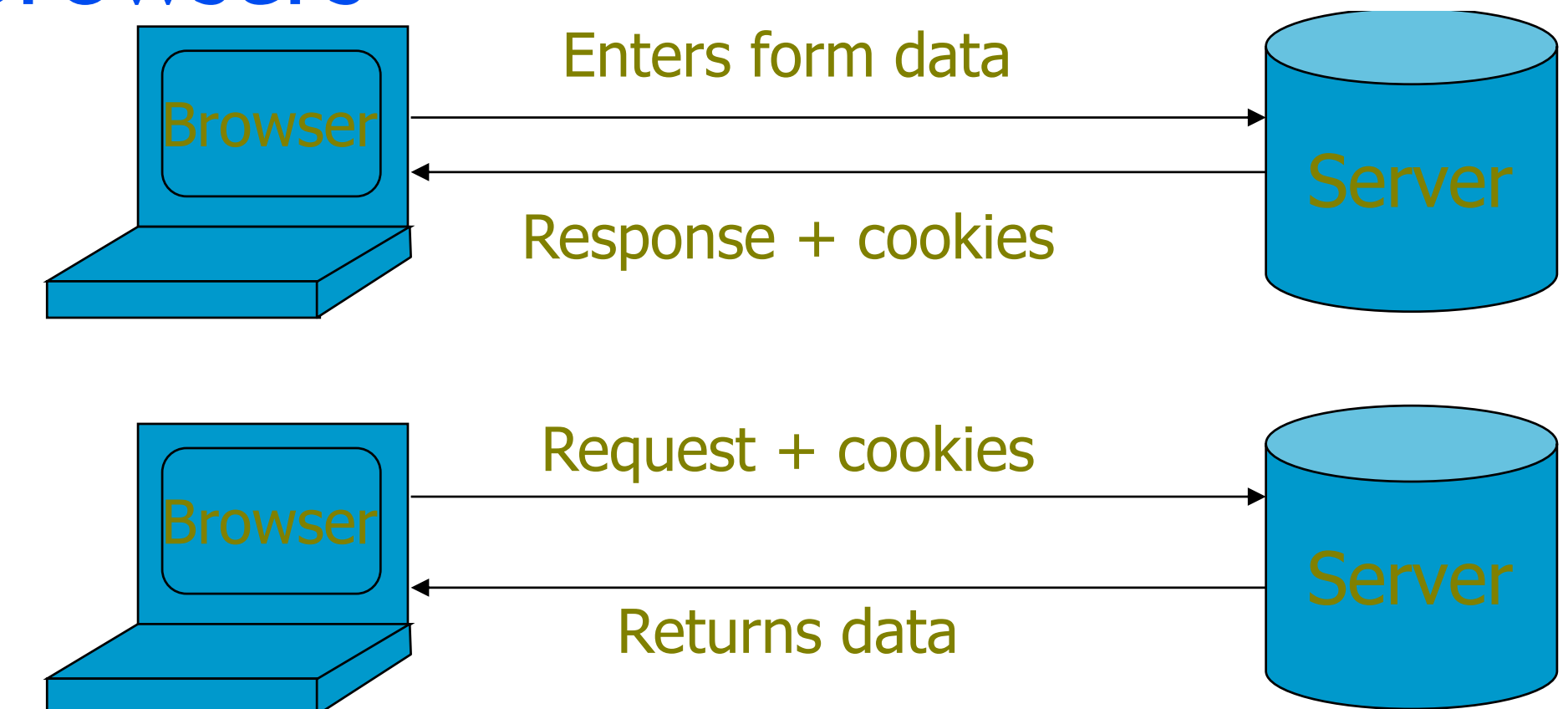
- Cookies were designed to offload server state to browsers

- ▶ Not initially part of web tools (Netscape)
- ▶ Allows users to have cohesive experience
- ▶ E.g., flow from page to page,

- Someone made a design choice

- ▶ Use cookies to *authenticate* and *authorize* users
- ▶ E.g. Amazon.com shopping cart, WSJ.com

- Q: What is the threat model?



## Cookies

A cookie is a name/value pair created by a website to store information on your computer





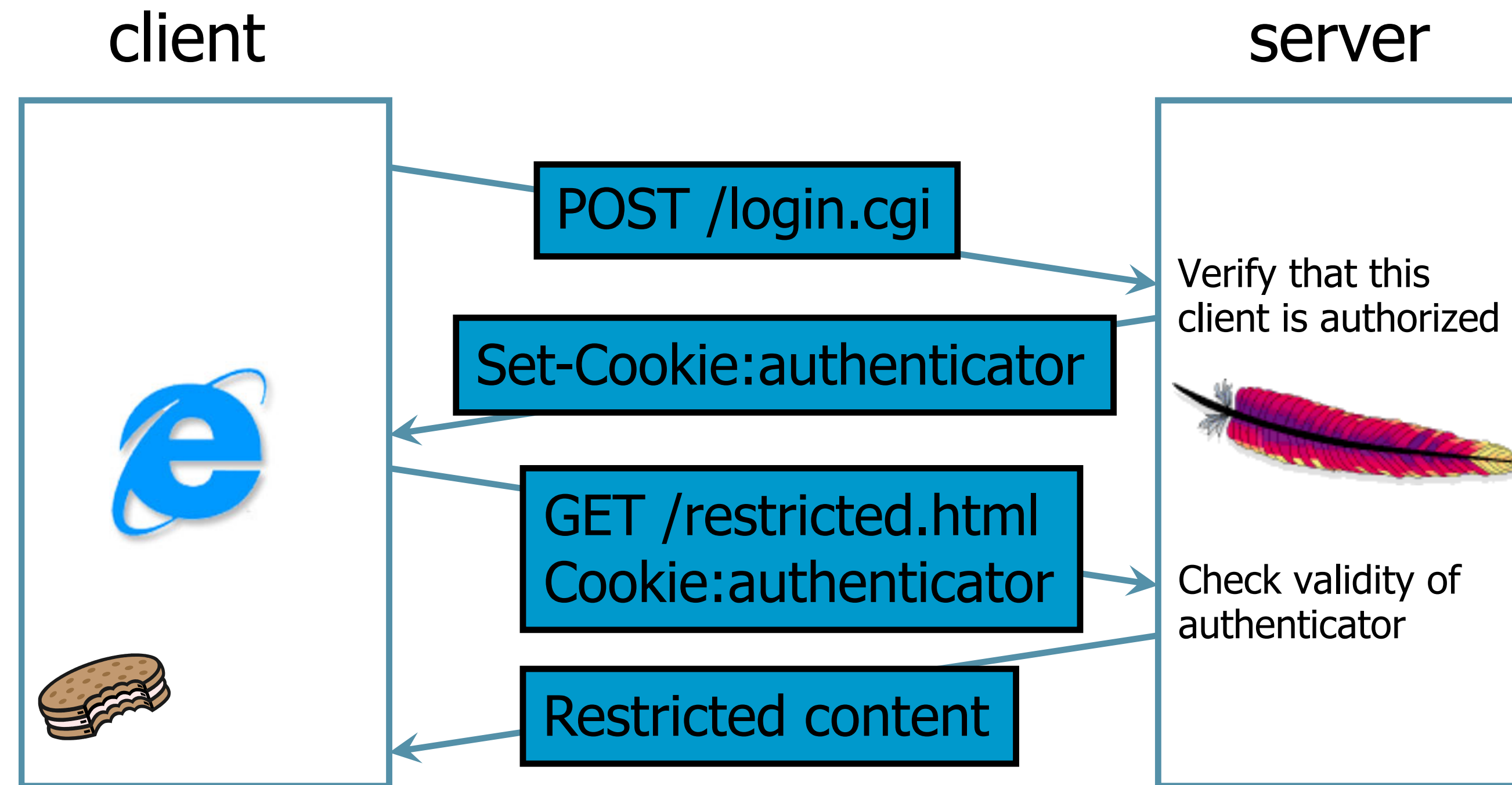
- An example cookie from my browser

Name	session-token
Content	"s7yZiOvFm4YymG...."
Domain	.amazon.com
Path	/
Send For	Any type of connection
Expires	Monday, September 08, 2031 7:19:41 PM

- Stored by the browser and used by the web applications
  - ▶ used for authenticating, tracking, and maintaining specific information about users
  - ▶ e.g., site preferences, contents of shopping carts
  - ▶ data may be sensitive
  - ▶ may be used to gather information about specific users
- Cookie ownership: Once a cookie is saved on your computer, only the website that created the cookie can read it

- HTTP is stateless
  - ▶ How does the server recognize a user who has signed in?
- Servers can use cookies to store state on client
  - ▶ After client successfully authenticates, server computes an authenticator and gives it to browser in a cookie
    - Client cannot forge authenticator on his own (session id)
  - ▶ With each request, browser presents the cookie
  - ▶ Server verifies the authenticator
  - ▶

# A Typical Session with Cookies



Authenticators must be unforgeable and tamper-proof  
(malicious clients shouldn't be able to modify an existing authenticator)

**How to design it?**

- **New design choice means**
  - ▶ Cookies must be protected
    - Against forgery (integrity)
    - Against disclosure (confidentiality)
- **Cookies not robust against web designer mistakes, committed attackers**
  - ▶ Were never intended to be
  - ▶ Need the same scrutiny as any other tech.



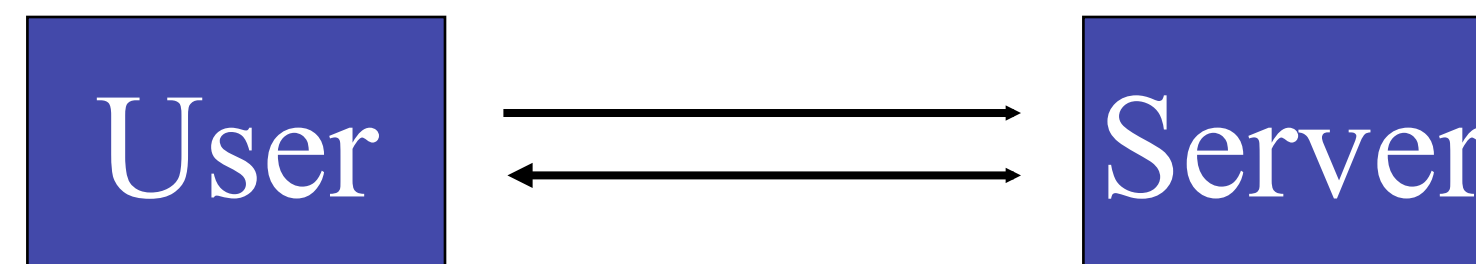
**Many security problems arise out of a technology built for one thing incorrectly applied to something else.**

# Cookie Design 1: mygorilla.com

- Requirement: authenticate users on site

myschool.com

- Design:
  1. set cookie containing hashed username
  2. check cookie for hashed username



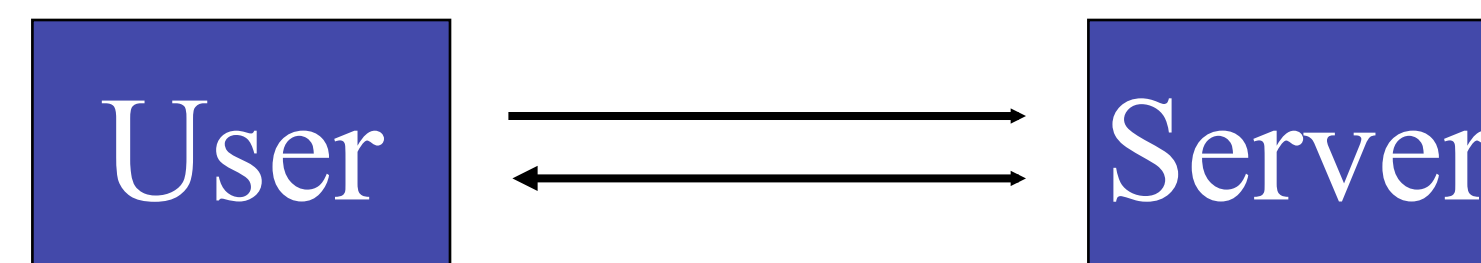
- Q: Is there anything wrong with this design?

# Cookie Design 2: mygorilla.com

- Requirement: authenticate users on site

myschool.com

- Design:
  1. set cookie containing **encrypted** username
  2. check cookie for **encrypted** username



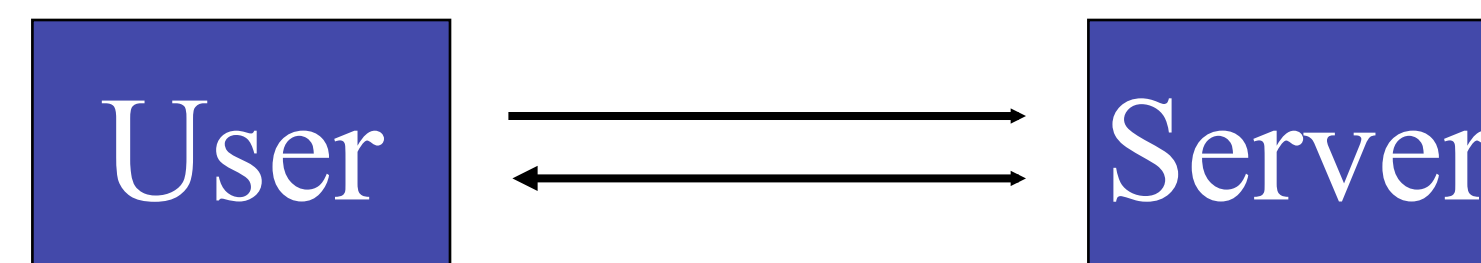
- Q: Is there anything wrong with this design?

# Cookie Design 2: mygorilla.com

- Requirement: authenticate users on site

## myschool.com

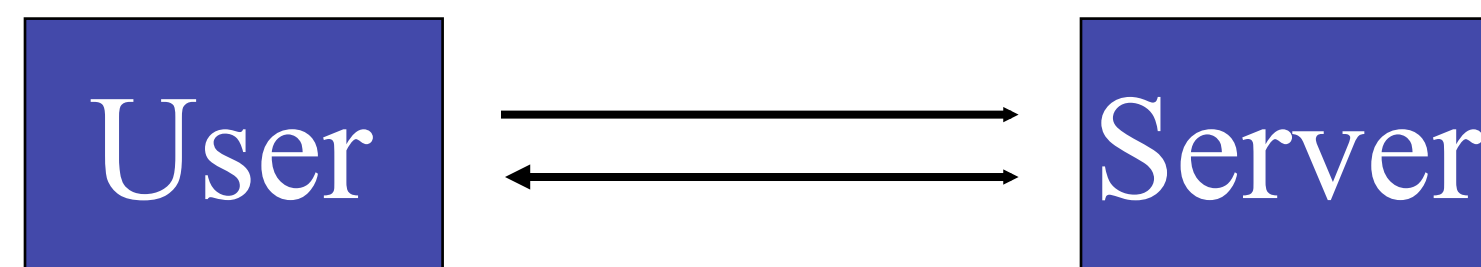
- Design:
  1. set cookie containing **encrypted** + **HMAC'd** username
  2. check cookie for **encrypted** + **HMAC'd** username



- Q: Is there anything wrong with this design?

# Exercise: Cookie Design

- Design a secure cookie for **myschool.com** that meets the following requirements
- Requirements
  - ▶ Users must be authenticated (assume digest completed)
  - ▶ Time limited (to 24 hours)
  - ▶ Unforgeable (only server can create)
  - ▶ Privacy-protected (username not exposed)
  - ▶ Location safe (cannot be replayed by another host)

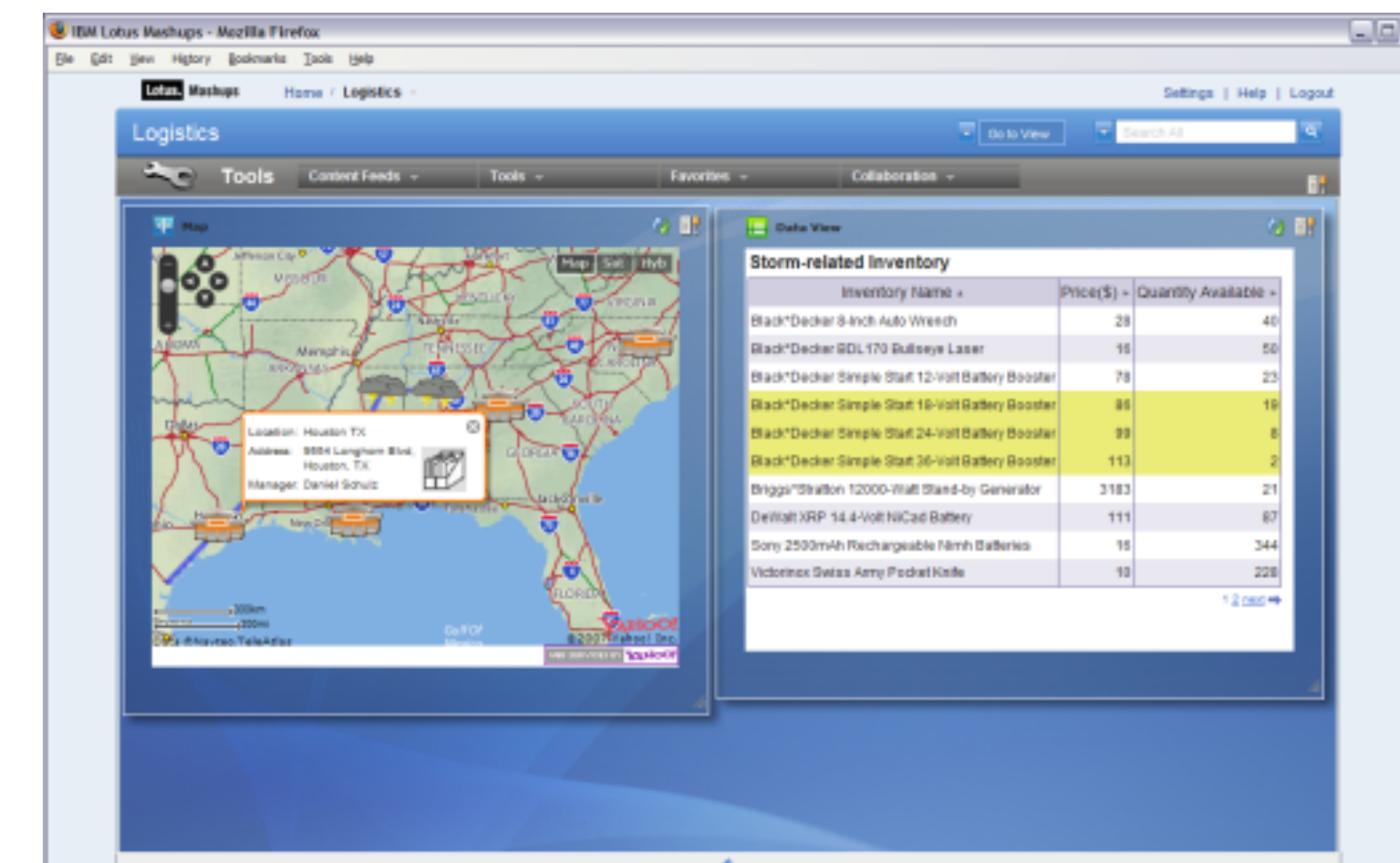


$E\{k_s, "host\_ip : timestamp : username"\} + HMAC\{k_s, "..."\}$



# Content from Multiple Sites

- Browser stores cookies from multiple websites
  - ▶ Tabs, mashups, ...
- **Q. What is the threat model?**
- More generally, browser stores *content* from multiple websites
  - ▶ HTML pages
  - ▶ Cookies
  - ▶ Flash
  - ▶ Java applets
  - ▶ JavaScript
- How do we isolate content from multiple sites?



# Client Side Scripting

- Web pages (HTML) can embed dynamic contents (code) that can be executed on the browser
- JavaScript
  - ▶ embedded in web pages and executed inside browser
- Java applets
  - ▶ small pieces of Java bytecodes executed in browsers
  - ▶

```
<html>
  ...
  <P>
<script>
  var num1, num2, sum
  num1 = prompt("Enter first number")
  num2 = prompt("Enter second number")
  sum = parseInt(num1) + parseInt(num2)
  alert("Sum = " + sum)
</script>
  • ...
  • </html>
```

Browser receives content, displays HTML and executes scripts

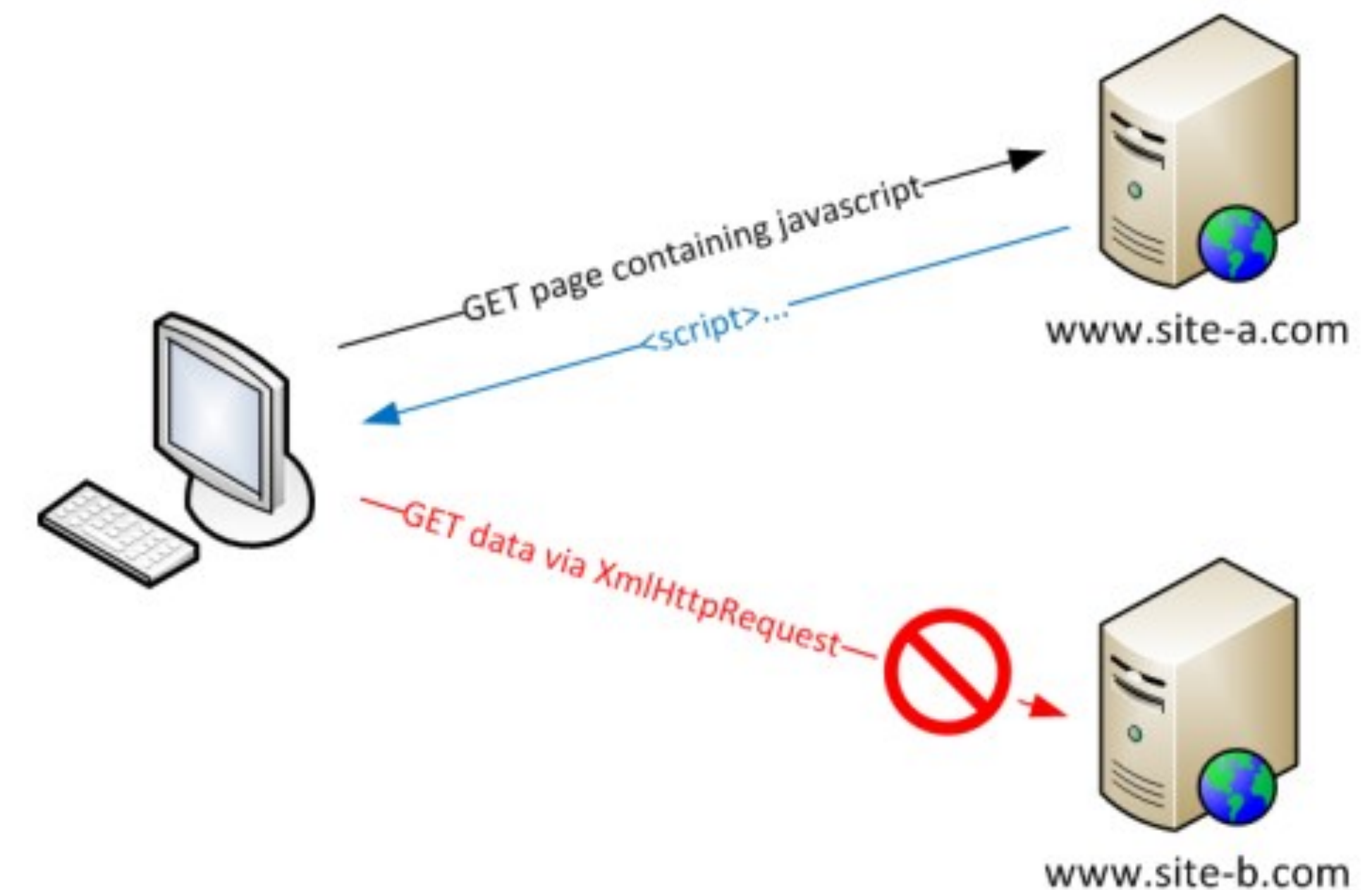
Client-side scripting can access (read/write) the following resources

- Local files on the client-side host
- Webpage resources maintained by the browser: Cookies, Domain Object Model (DOM) objects
  - steal private information
  - control what users see
  - impersonate the user

- Web users visit multiple websites simultaneously
- A browser serves web pages (which may contain programs) from different web domains
  - ▶ i.e., a browser runs programs provided by mutually untrusted entities
  - ▶ Running code one does not know/trust is dangerous
  - ▶ A browser also maintains resources created/updated by web domains
- Browser must confine (sandbox) these scripts so that they cannot access arbitrary local resources
- Browser must have a security policy to manage/protect browser-maintained resources and to provide separation among mutually untrusted scripts

# Same-Origin Policy

- A set of policies for isolating content (scripts and resources) across different sites (*origins*)
  - ▶ E.g., evil.org scripts cannot access bank.com resources.
- What is an origin?
  - ▶ site1.com vs site2.com?
    - Different hosts are different origins
  - ▶ http://site.com vs https://site.com?
    - Different protocols are different origins
  - ▶ http://site.com:80 vs http://site.com:8080?
    - Different ports are different origins
  - ▶ http://site1.com vs http://a.site1.com?
    - Establishes a hierarchy of origins



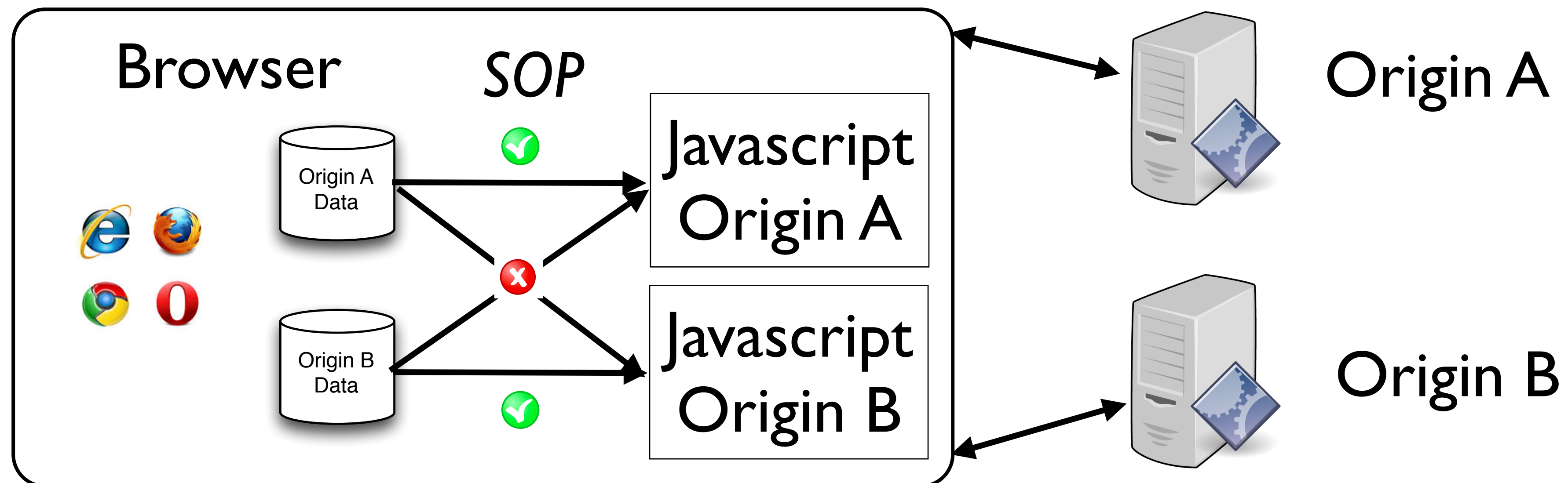


# SOP: What it Controls?

- Same-origin policy applies to the following accesses:
  - ▶ manipulating browser windows
  - ▶ URLs requested via the XMLHttpRequest
    - XMLHttpRequest is an API that can be used by web browser scripting languages to transfer XML and other text data to and from a web server using HTTP, by establishing an independent and asynchronous communication channel.
    - used by AJAX
  - ▶ manipulating frames (including inline frames)
  - ▶ manipulating documents (included using the object tag)
  - ▶ manipulating cookies
  - ▶

# Same-Origin Policy

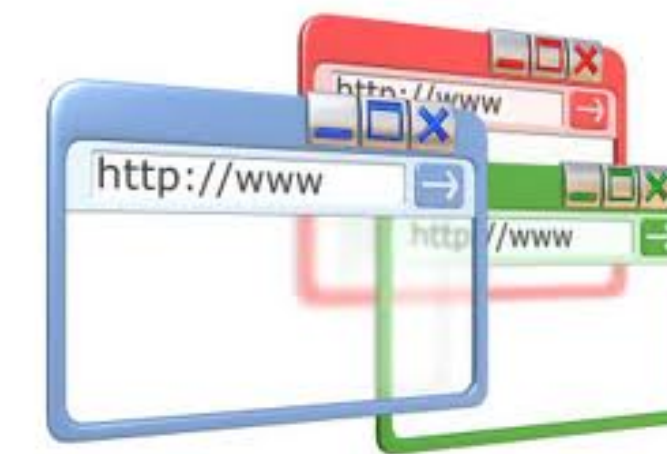
- *Principle:* Any active code from an origin can read only information stored in the browser that is from the same origin
  - ▶ Active code: Javascript, VBScript, ...
  - ▶ Information: cookies, HTML responses, ...





- Scripts from two origins in the same domain may wish to interact
  - ▶ www.example.com and program.example.com
- Any web page may set *document.domain* to a
  - ▶ “right-hand, fully-qualified fragment of its current host name” (example.com, but not ample.com)
- Then, **all scripts** in that domain may share access
  - ▶ All or nothing
- NOTE: Applies “null” for port, so does not actually share with normal example.com:80

- Complete and partial bypasses exist
  - ▶ Browser bugs
  - ▶ Limitations if site hosts unrelated pages
    - Example: Web server often hosts sites for unrelated parties
    - `http://www.example.com/account/`
    - `http://www.example.com/otheraccount/`
    - Same-origin policy allows script on one page to access document properties from another
  - ▶ Functionality often requires SOP bypass!
    - Many advertisement companies hire people to find and exploit SOP browser bugs for cross-domain communication
    - E.g., JSON with padding (JSONP)
- Cross-site scripting
  - ▶ Execute scripts from one origin in the context of another



# Cross Site Scripting (XSS)

- Recall the basics
  - ▶ scripts embedded in web pages run in browsers
  - ▶ scripts can access cookies
    - get private information
  - ▶ and manipulate DOM objects
    - controls what users see
  - ▶ scripts controlled by the same-origin policy
- Why would XSS occur
  - ▶ Web applications often take user inputs and use them as part of webpage

- Assume the following is posted to a message board on your favorite website which will be displayed to everyone:

**Hello message board.**

**<SCRIPT>malicious code</SCRIPT>**

**This is the end of my message.**

- Now a reasonable ASP (or some other dynamic content generator) uses the input to create a webpage (e.g., blogger nonsense).
- Anyone who view the post on the webpage can have local authentication cookies stolen.
- Now a malicious script is running
  - ▶ Applet, ActiveX control, JavaScript...



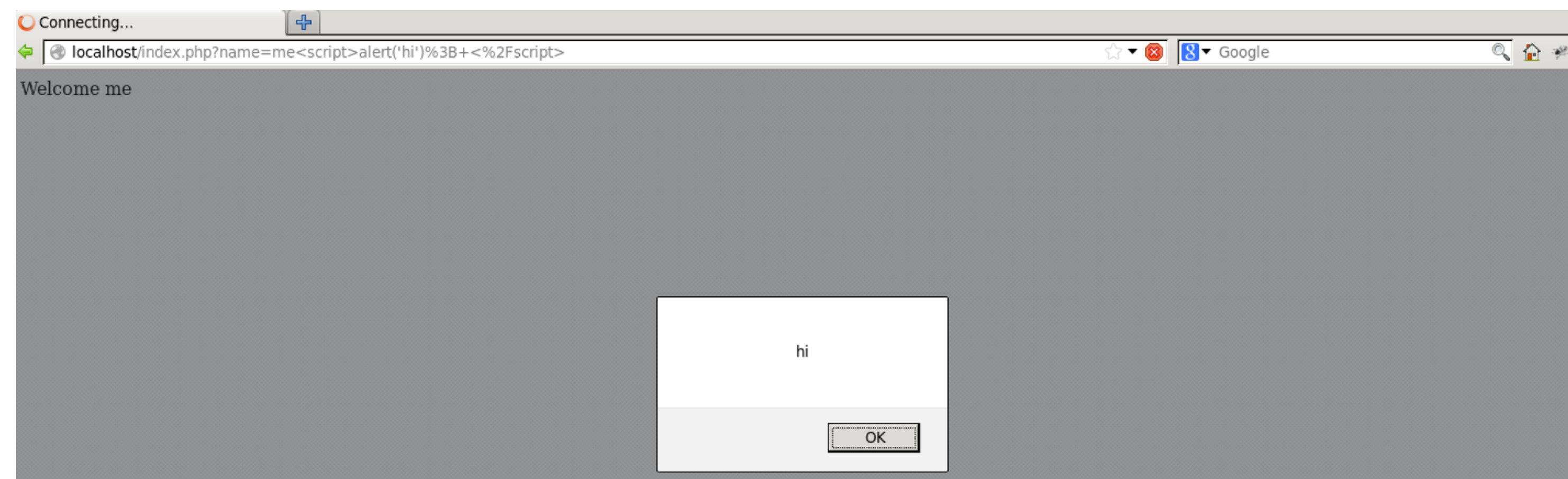
- Script from attacker is executed in the victim origin's context
  - ▶ Enabled by inadequate filtering on server-side
- Effects of Cross-Site Scripting
  - ▶ Can manipulate any DOM component on victim.com
  - ▶ Control links on page
  - ▶ Control form fields (e.g. password field) on this page and linked pages.
  - ▶ Can infect other users: MySpace.com worm
- Three types
  - ▶ Reflected
  - ▶ Stored
  - ▶ DOM Injection



```
<?php
$name = $_GET['name'];
echo "Welcome $name<br>";
?>

<form method="get" action="index.php">
  Name: <input type="text" name="name" /><br />
  <input type="submit" value="submit" />
</form>
```

`index.php?name=guest<script>alert('hi')</script>`



# MySpace.com (Samy worm)

- Users can post HTML on their pages
  - ▶ MySpace.com ensures HTML contains no `<script>`, `<body>`, `onclick`, `<a href=javascript://>`
  - ▶ However, attacker find out that a way to include Javascript within CSS tags:  
`<div style="background:url('javascript:alert(1)')">`
  - ▶ And can hide “javascript” as “java\nscript”
- With careful javascript hacking:
  - ▶ Samy’s worm: infects anyone who visits an infected MySpace page ... and adds Samy as a friend.
  - ▶ Samy had millions of friends within 24 hours.
- More info: <http://namb.la/popular/tech.html>

# Web Systems Evolve ...

- The web has evolved from a *document retrieval* and rendering to sophisticated *distributed application platform* providing:

- ▶ dynamic content
- ▶ user-driven content
- ▶ interactive interfaces
- ▶ multi-site content
- ▶ ....



- With new interfaces comes new vulnerabilities ...



# Cross-site Request Forgery

- An XSS attack exploits the trust the browser has in the server to filter input properly
- A CSRF attack exploits the trust the server has in a browser

- ▶ Authorized user submits unintended request

- Attacker Maria notices weak bank URL `GET http://bank.com/transfer.do?acct=BOB&amount=100 HTTP/1.1`

- Crafts a malicious URL `http://bank.com/transfer.do?acct=MARIA&amount=100000`

- Exploits social engineering to get Bob to click the URL

```
<a href="http://bank.com/transfer.do?acct=MARIA&amount=100000">View my Pictures!</a>
```

- Can make attacks not obvious

```

```

- ▶ Defense: Referrer header

- Bank does not accept request unless referred to (linked from) the bank's own webpage
- Disadvantage: privacy issues

- More Example:
  - ▶ User logs in to bank.com. Forgets to sign off.
  - ▶ Session cookie remains in browser state
- Then user visits another site containing:

```
<form name=F action=http://bank.com/BillPay.php>  
<input name=recipient value=badguy> ...  
<script> document.F.submit(); </script>
```

  - ▶ Browser sends user auth cookie with request
  - ▶ Transaction will be fulfilled
- Problem: The browser is a **confused deputy**; it is serving both the websites and the user and gets confused who initiated a request
- [https://www.youtube.com/watch?v=5joXlSkQtVE&feature=emb\\_logo](https://www.youtube.com/watch?v=5joXlSkQtVE&feature=emb_logo)

# HTTP Response Splitting

- Again, due to insufficient server-side filtering
  - ▶ Cookies can be set to arbitrary values to split HTTP response

```
String author = request.getParameter(AUTHOR_PARAM);  
...  
Cookie cookie = new Cookie("author", author);  
cookie.setMaxAge(cookieExpiration);  
response.addCookie(cookie);
```

```
HTTP/1.1 200 OK  
...  
Set-Cookie: author=Jane Smith  
...
```

```
HTTP/1.1 200 OK  
...  
Set-Cookie: author=Wiley Hacker  
  
HTTP/1.1 200 OK  
...
```

- ▶ Can be used for page hijacking through proxy server

- Virtual sessions are implemented in many ways
  - ▶ session ID in cookies, URLs
  - ▶ If I can *guess*, *infer*, or *steal* the session ID, game over
  - ▶ Login page using HTTPS, but subsequent communication is not! Cookies sent in cleartext
  - ▶ If your bank encodes the session ID in the url, then a malicious attacker can simply keep trying session IDs until gets a good one.
  - ▶ ... note that if the user was logged in, then the attacker has full control over that account. <http://www.mybank.com/loggedin?sessionid=11>
  - ▶ Countermeasure: HTTPS, secure cookie design

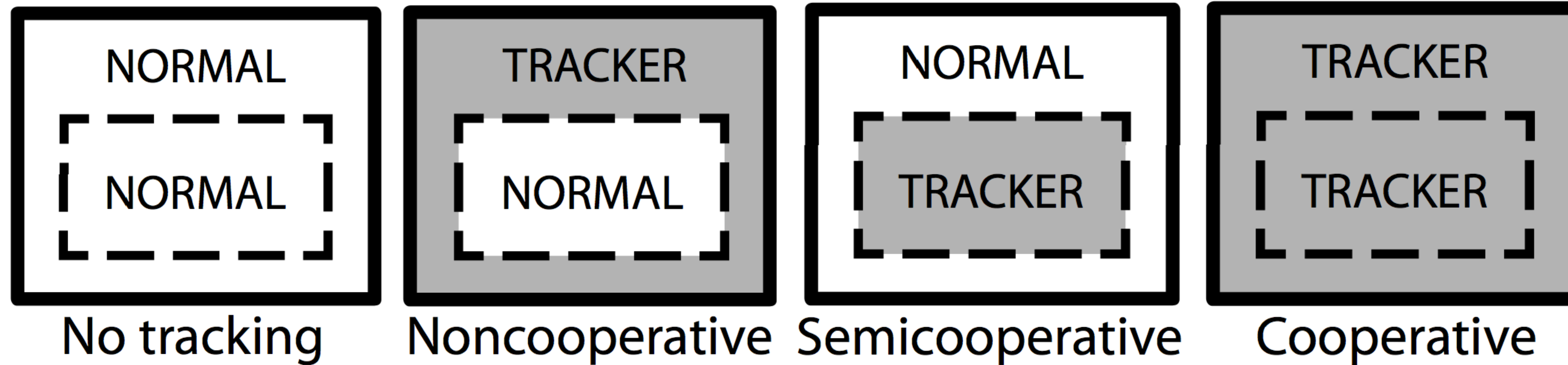
- Have you ever ...
  - ▶ Searched for a product on some website
  - ▶ ...Advertisement for the same product shows up on another website?
  - ▶ **Reason:** Tracking! Profile users for targeted advertisement
- Study by WSJ found (2012)
  - ▶ 75% of top 1000 sites feature social networking plugins
    - Can match users' identities with web-browsing activities
- abine and UC Berkeley found
  - ▶ Online tracking is 25% of browser traffic
    - 20.28% google analytics
    - 18.84% facebook



<http://www.abine.com/>

- Tracking is done in following configurations

Protecting Browser State from Web Privacy Attacks : Jackson et al.



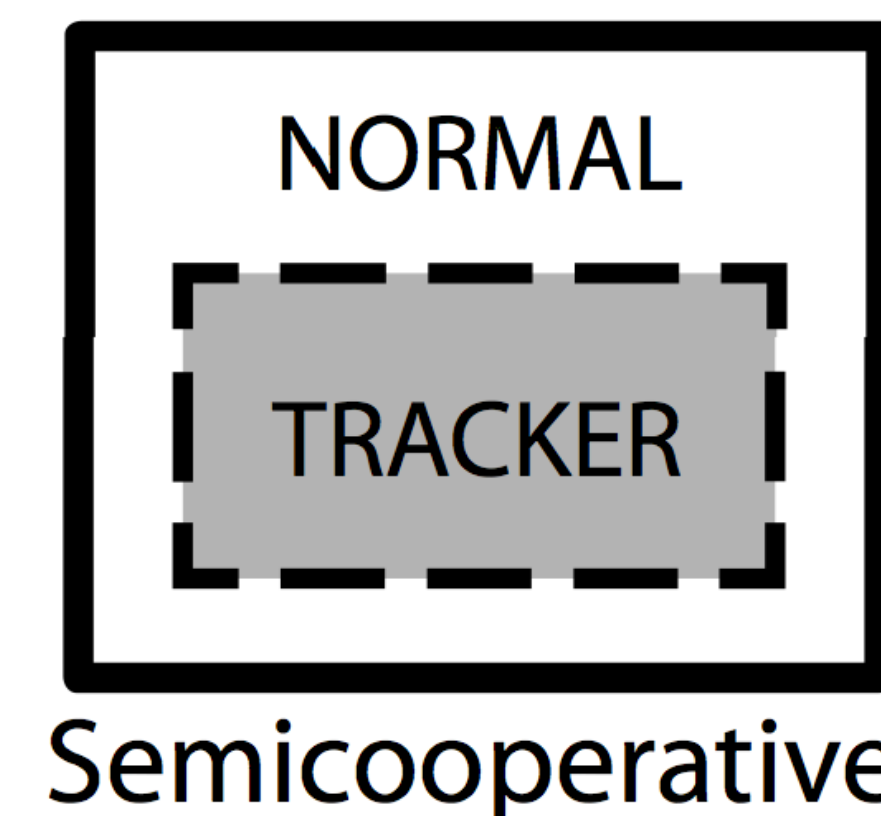
- “Tracker” code is from
  - Social networking sites
  - Analytics
  - Advertisement agencies
  - ...

- Objective of tracking code is to maintain state of users across multiple sites
  - ▶ Build profile of sites visited
- Semi-cooperative tracking done by
  - ▶ Javascript
    - e.g., Cached redirect URLs
  - ▶ Web bugs
    - IxI images
    - Ever wondered why email clients have “Display images”?
  - ▶ IFrames
  - ▶ Cookies
    - Traditional, flash, HTML5 LocalStorage, ...
- **Tasks:** (1) get your tracking code running; (2) store state; (3) send to server



# Third-Party Cookies

- A third-party cookie is a cookie from a website different from the website being viewed
- Browsers can block third-party cookies
  - ▶ Different browsers have different variations
    - Some completely block
    - “Do Not Track” - except Chrome
- Limitation
  - ▶ Other ways exist to store state (more)
    - Canvas fingerprinting
    - Evercookies
    - “Cookie syncing”
- OpenWPM - <https://github.com/citp/OpenWPM>





# Unintended Tracking

- “Data” access not all governed by same-origin policy

- ▶ Specified: HTML DOM, cookies

- ▶ What about

- Web caches?

- ▶ Tracking notes time to fetch URL

- ▶ If URL in cache, served faster

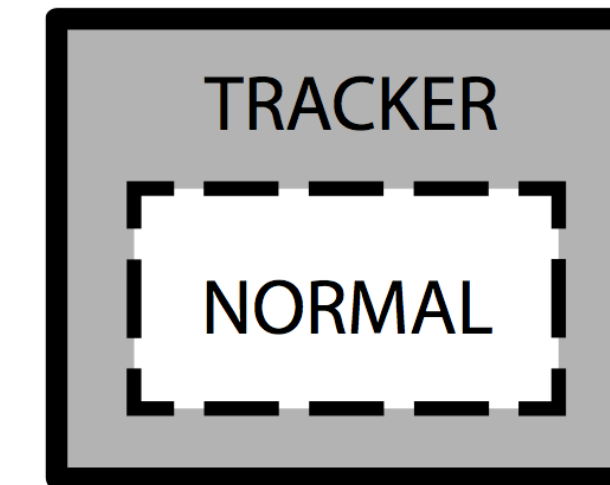
- Visited links?

- ▶ Mostly fixed in current browsers

- **Take-away:** Difficult to prevent tracking if *any* browser state is stored

- To mitigate tracking

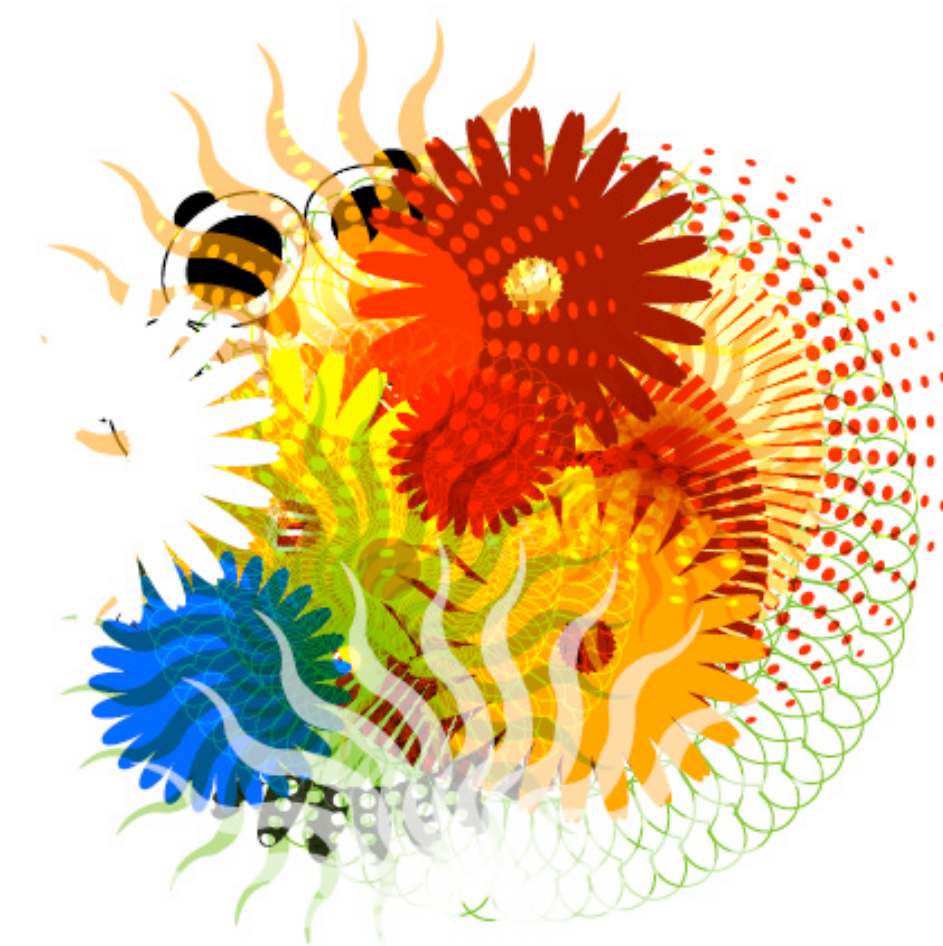
- ▶ Reset browser regularly, store no state, visit random sites!



```
a { color: blue; }
a:visited { color: red; }

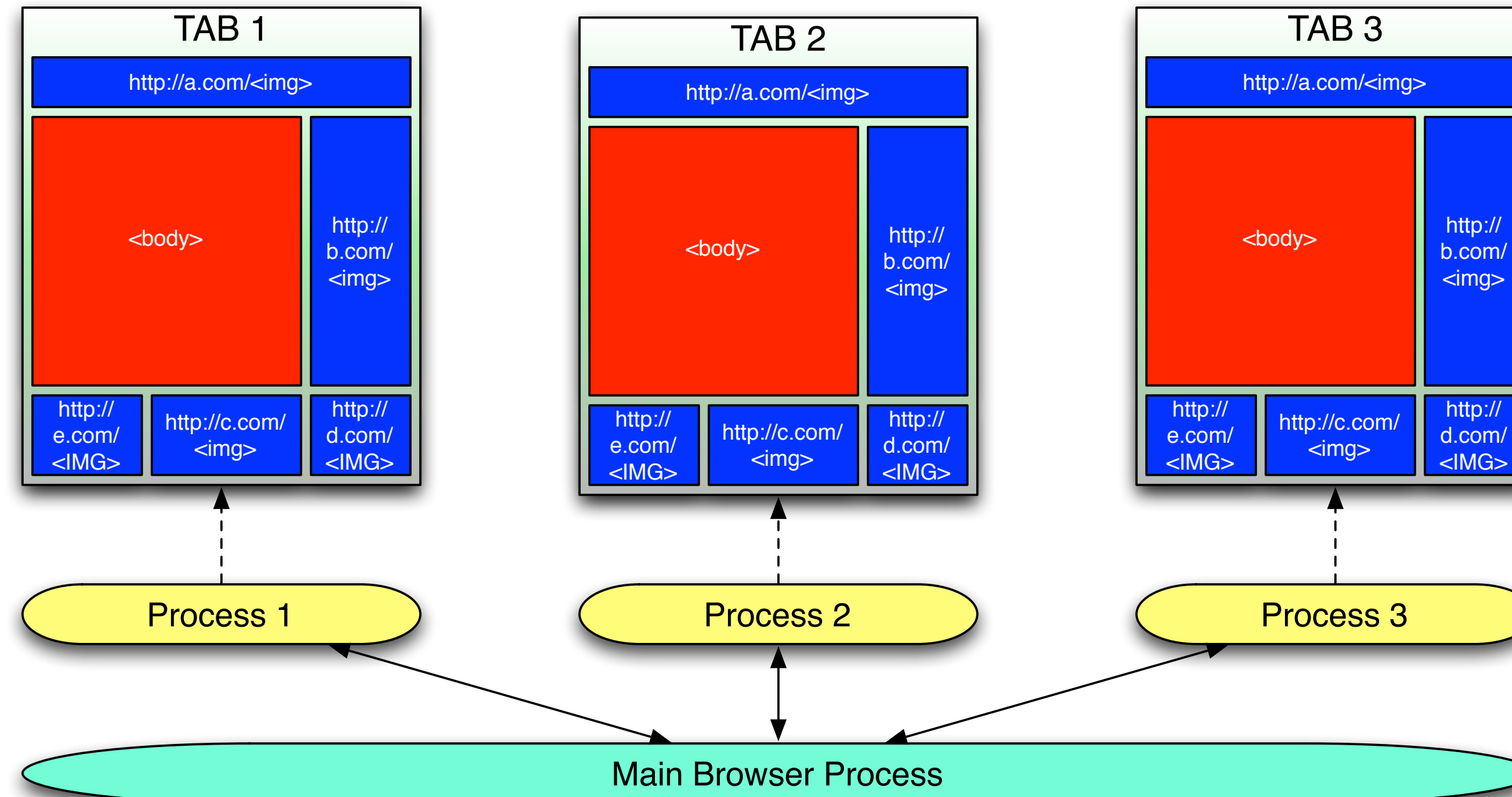
if (document.getElementById('jones').currentStyle.color=='red')
  document.writeln('<p>Hello! I see you\'ve been to Jones.</p>');
  document.writeln('Don\'t buy from Jones - their widgets');
  document.writeln('are made from recycled babies.</p>');
```

- Browsers are the new operating systems
- Huge, complex systems that support
  - ▶ Many document types, structures, e.g., HTML, XML, ...
  - ▶ Complex rendering, e.g., CSS, CSS 2.0
  - ▶ Many “program/scripting” languages, e.g., JavaScript
  - ▶ Dynamic content, e.g., AJAX
  - ▶ Native code execution, e.g., ActiveX



- Virtualized computers in a single program ...

- We don't have the ability to control this much complexity, so we have to try other things ...
  - ▶ Restricting functionality, e.g., NoScript
  - ▶ Process Isolation, e.g., OP, Chrome
    - Read: <http://www.google.com/googlebooks/chrome/>



- What did they do to build a more secure browser?
- (I) Decompose the browser into multiple processes
  - Called “Privilege Separation”
- What are the permissions of a set of processes forked from the same parent?

- What did they do to build a more secure browser?
- (1) Decompose the browser into multiple processes
  - Called “Privilege Separation”
- What are the permissions of a set of processes forked from the same parent?  
Same as parent
- (2) Need different policy for each process
  - **Multiple subjects** in the access control policy
- What browser processes are trusted to manage the permissions?

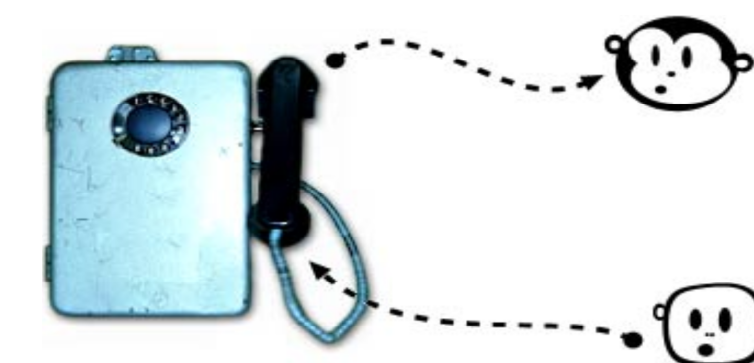
- What did they do to build a more secure browser?
- (1) Decompose the browser into multiple processes
  - Called “Privilege Separation”
- What are the permissions of a set of processes forked from the same parent?  
Same as parent
- (2) Need different policy for each process
  - Multiple subjects in the access control policy
- What browser processes are trusted to manage the permissions? None
- (3) Need mandatory access control
  - Subjects cannot escape confined “**protection domain**”

- How do you determine what parts of the browser should be a “subject” and identify the permissions to be assigned to that subject?
- One subject (client)
  - Code that requires the same permissions to run
  - E.g., a particular web page
- Another subject (**server**)
  - Code that manages the same permissions
  - E.g., UI, network, and storage subsystems
- How do we determine the **permission assignments**?

- How do you determine what parts of the browser should be a “subject” and identify the permissions to be assigned to that subject?
- One subject (client)
  - Code that requires the same permissions to run
  - E.g., a particular web page
- Another subject (**server**)
  - Code that manages the same permissions
  - E.g., UI, network, and storage subsystems
- How do we determine the **permission assignments**?
  - Least privilege
  - Information flow

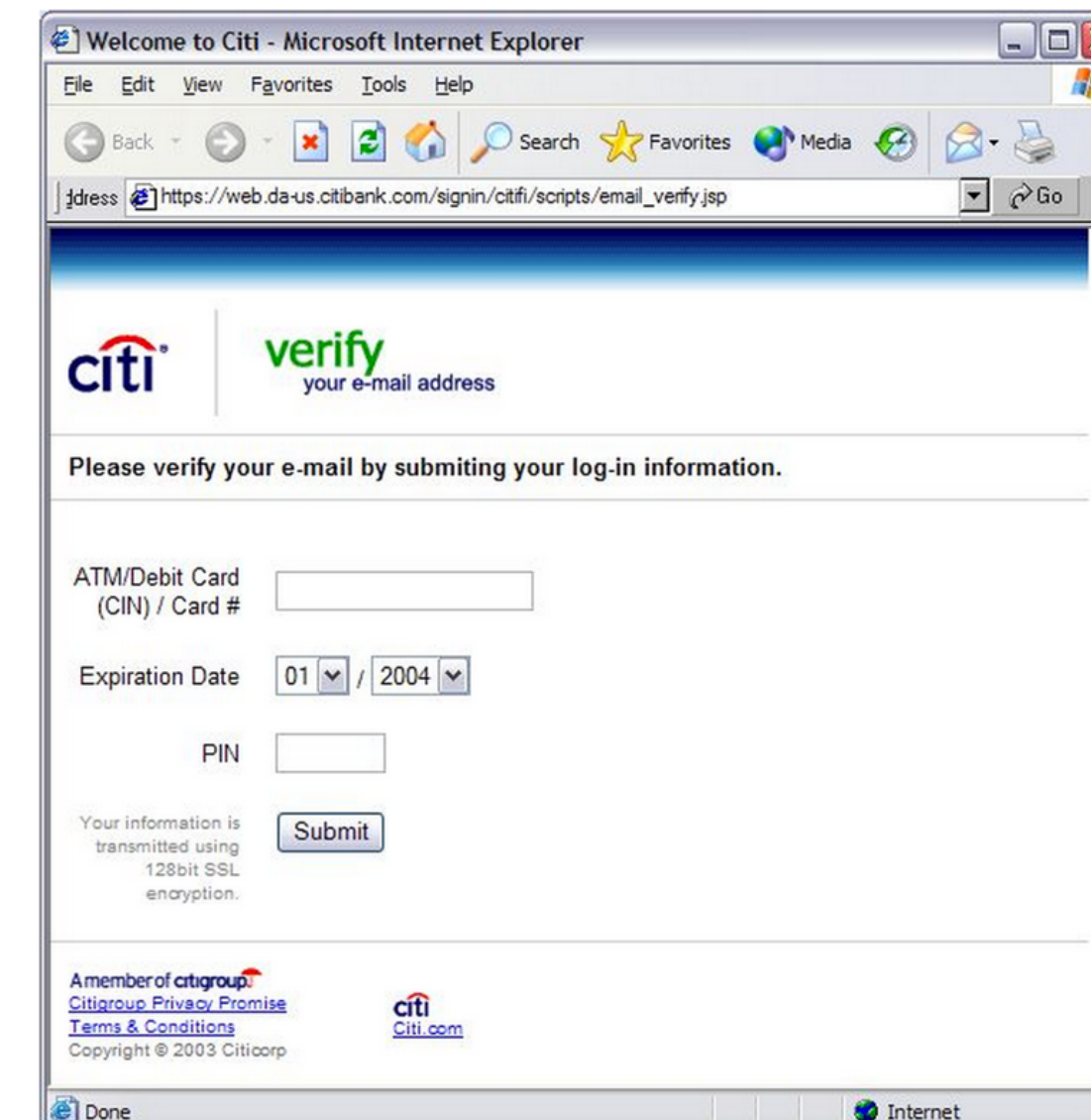
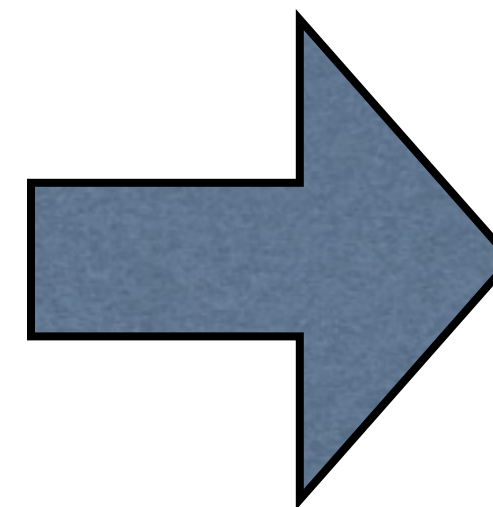
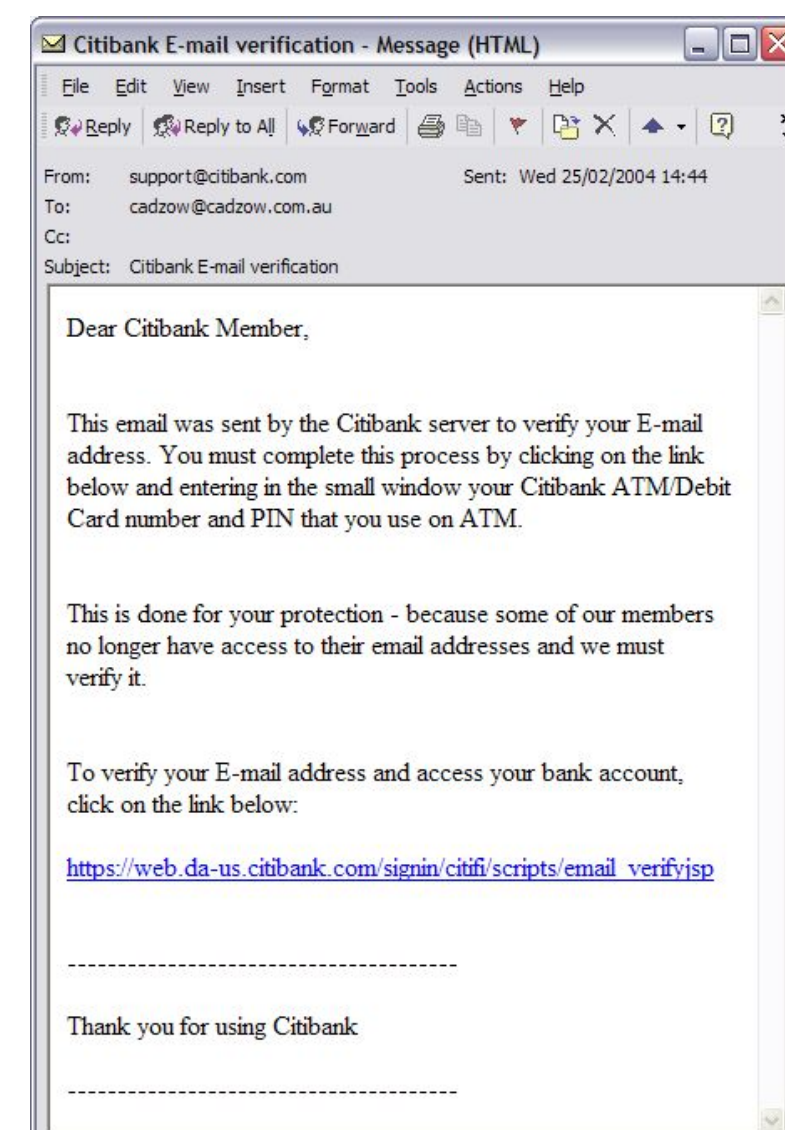


- A *plugin* is simply a program used by a browser to process content
  - ▶ MIME type maps content to plugin
  - ▶ Like any old application (e.g., RealAudio)
  - ▶ Newer browsers have autoinstall features
- Plugins are sandboxed, but have been circumvented in various ways
  - ▶ Interesting design point - Google Chrome allows “native” plugins but still preserves (some) security!
    - Native Client sandbox for running compiled C/C++ code
- **Moral:** beware of plugins



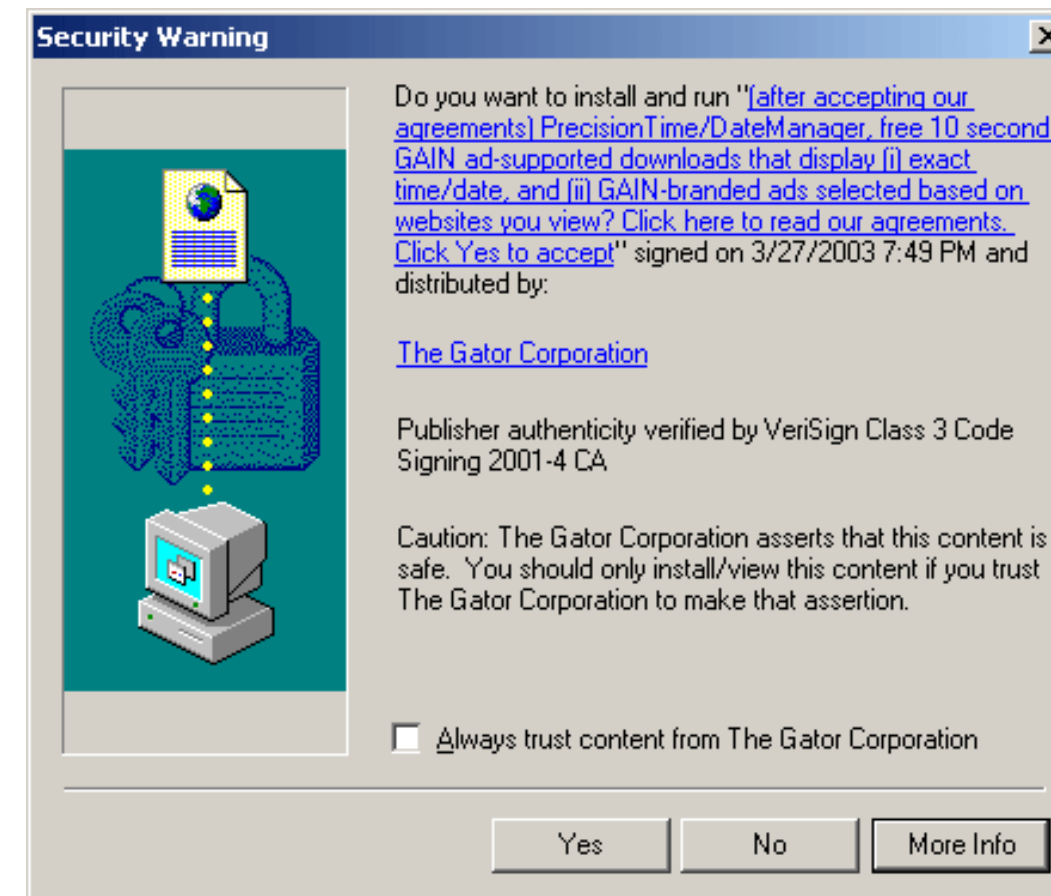
# Social Engineering

- Attacks another weak point -- users!
- **Phishing**
  - ▶ Lure users using bait (fishing) to steal valuable information
  - ▶ Common technique: mimic original site and use similar URL
    - www.aol.com vs www.aol.com
    - Combine with other techniques e.g., turn off address bar



# Drive by downloads

- Using a deceptive means to get someone to install something on their own (spyware/adware)



- ▶ Often appears as an error message on the browser
- ▶ Sometimes, user does not click anything at all!
- ▶ **Concern:** *extortion-ware* -- pay us \$ to unencrypt your data
  - Used to demand \$ for uninstall of annoying software
- ▶ “biggest cybersecurity threat” - Kaspersky
- **Answer:** Back up stuff externally that you really want!

- An injection that exploits the fact that many inputs to web applications are
  - ▶ under control of the user
  - ▶ used directly in SQL queries against back-end databases
- Bad form inserts escaped code into the input ...

```
xUserId = getRequestString("UserId");  
txtSQL = "SELECT * FROM Users WHERE UserId = " + xUserId;
```

- This vulnerability became one of the most widely exploited and costly in web history.
  - ▶ Industry reported as many as 16% of websites were vulnerable to SQL injection in 2007
  - ▶ This may be inflated, but has been an ongoing problem.

- An injection that exploits the fact that many inputs to web applications are
  - ▶ under control of the user
  - ▶ used directly in SQL queries against back-end databases
- Bad form inserts escaped code into the input ...

```
SELECT email, login, last_name
FROM user_table
WHERE email = 'x'; DROP TABLE members; --';
```

- This vulnerability became one of the most widely exploited and costly in web history.
  - ▶ Industry reported as many as 16% of websites were vulnerable to SQL injection in 2007
  - ▶ This may be inflated, but has been an ongoing problem.

- Largely just applications
  - In as much as application are secure
  - Command shells, interpreters, are dangerous
- Broad Approaches
  - Validate input (also called *input sanitization*)
  - Limit program functionality
    - Don't leave open ended-functionality
  - Execute with limited privileges
  - Input tracking, e.g., *taint tracking*
  - Source code analysis, e.g., c-cured



- Web security has to consider threat models involving several parties
  - Web browsers
  - Web servers
  - Web applications
  - Users
  - Third-party sites
  - Other users
- Security is so difficult in the web because it was largely *retrofitted*
- *zzz*

