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CSE543 - Computer Security

**CSE543** Introduction to Computer and Network Security Module: Authentication



## Authentication and Authorization

- Fundamental mechanisms to enforce security on a system
- Authentication: Identify the principal responsible for a "message"
  - Distinguish friend from foe
- the identity of a principal
  - restricted operation
- Today, we discuss principles behind authentication





## Authorization: Control access to system resources based on

Determine whether a principal has the permissions to perform a









## What is Authentication?

- Short answer: establishes identity
  - Answers the question: To whom am I speaking?
- credentials
  - Credential is proof of identity
  - between credential and claimed identity
    - for some purpose
    - under some policy (what constitutes a good cred.?)



# Long answer: evaluates the authenticity of identity by proving

Evaluation – process that assesses the correctness of the association





## Why authentication?

- Well, we live in a world of rights, permissions, and duties
  - rights

## • Q: How does this relate to security?



## Authentication establishes our identity so that we can obtain the set of

 E.g., we establish our identity with Tiffany's by providing a valid credit card which gives us rights to purchase goods ~ physical authentication system

TIFFANY&CO.







## Why authentication (cont.)?

- Same in online world, just different constraints
  - of providing identity
    - e.g., by providing credit card *number* ~ electronic authentication system
  - Risks (for customer and vendor) are different
    - Q: How so?

management, and application of authentication systems.





## Vendor/customer are not physically co-located, so we must find other ways

## Computer security is crucially dependent on the proper design,





## What is Identity?

- That which gives you access ... which is largely determined by context
  - We all have lots of identities
  - Pseudo-identities
- Really, determined by who is evaluating credential
  - Driver's License, Passport, SSN prove ...
  - Credit cards prove ...
  - Signature proves ...
  - Password proves ...
  - Voice proves ...
- Exercise: Give an example of bad mapping between ider for which it was used.





rpose





## Credentials

- ... are evidence used to prove identity
- Credentials can be
  - Something I am
  - Something I have
  - Something I know

## INTERNATIONAL THEOLOGICAL UNIVERSITY

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## ADMINISTRATIVE CREDENTIAL

## YOUR NAME GOES HERE

Granied this third day of January, year of our Lord, two thousand and two, in Pasadena, California, United States of America

Title Administrative CredentialMajors: Education AdministrationValid: 01-03-02 for Life Levels: Pre-School - Grade 12Special Training: Human Resource Management and CurricUntegrityHon Rev. Professor Dr. Chief Alexanders Swift Cargle Justice D.D., H.D., J.D. Theologian; Academician Russian International Acade Science of Information, Communication, Control in Engineering, Nature, S and Management of Technology; Full Professor - International Econ Diploma of Honors - Communication (Space Federat Russia) - Chancellor of the University Dr. Mary Brane Eagle, Ph. D.		Majors: Education Administration	Minors: Counseling
		Special Training: Human Resource Management and Curriculum Developmen	
		Justice, tional Academy of ng, Nature, Society space Federation of Life ACO1030	



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## Something you know ...

- Passport number, mothers maiden name, last 4 digits of your social security, credit card number
- Passwords and pass-phrases
  - Note: passwords have historically been pretty weak
    - Same bias with the context. E.g.?
    - Passwords used in more than one place
  - Not just because bad ones selected: If you can remember it, then a computer can be a computer can be a computer of the selected of the selected. guess it
    - Computers can often guess very quickly
    - Easy to mount offline attacks
    - Easy countermeasures for online attacks





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## "Hoist with his own petard"

- The rule of seven plus or minus two.
  - George Miller observed in 1956 that most humans can remember about 5-9 things more or less at once.
  - Thus is a kind of maximal entropy that one can hold in your head.
  - This limits the complexity of the passwords you can securely use, i.e., not write on a sheet of paper.
  - A perfectly random 8-char password has less entropy than a 56-bit key.
- Implication?







## Password Storage

- Store password as a "hash" of its value
- What properties must hash function satisfy for this purpose?
  - Should hash entries be invertible?
  - Could two passwords result in the same hash value?



## Password Storage

- Store password as a "hash" of its value
  - Originally stored in /etc/passwd file (readable by all)
  - Now in /etc/shadow (readable only be root)
- What if an adversary can gain access to a password file? How would you attack this?



## "Salt" ing passwords

- Suppose you want to avoid a offline dictionary attack
  - bad guy precomputing popular passwords and looking at the password file
- A salt is a random number added to the password differentiate passwords when stored in /etc/shadow

## consequence: guesses each password independently





## Password Cracking

- Attacker can access the hashed password
  - Can guess and test passwords offline
- Called "password cracking"
- Lots of help
  - John the Ripper
- How well do these work?





## Cracking Passwords

- How hard are passwords to crack?
- How many 8-character passwords available?



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## • How many 8-character passwords are there given that 128 characters are

## Cracking Passwords

- How hard are passwords to crack?
- - $128^8 = 2^{56}$
- How many guesses to find one specific user's password?
  - $2^{56}/2 = 2^{55}$



## • How many 8-character passwords given that 128 characters are available?

## Guess Again...

- How do you know if your password will be guessed?
  - Follow password-composition policies
- Example properties
  - Length: 8 or 12 or 16 chars?
  - **Requirements**: Password must contain at least one...
  - Blacklist: Password must not contain a dictionary word
- How do you know which policy to choose?
  - Studied in "Guess again ...: Measuring password strength by simulating password cracking algorithms," Gage Kelley, et al., IEEE Security and Privacy, 2012



## Guess Number

- How do you predict how many guesses it will take to crack your password?
  - Try to crack it?
    - That can be time consuming
  - Compute number of guesses it would take?
    - How do we do that?





## Guess Number

- to crack a specific password
  - Produce a deterministic guess ordering
- For "brute-force Markov" cracker
  - Uses frequencies of start chars and following chars
    - Most likely first, most likely to follow that, and so on...
  - Sum the number of guesses to find each character
    - In an N character alphabet and a password of length L:
      - ► The first character is the kth char tried in (k-1)N<sup>L-1</sup> guesses
      - The second character is the kth char tried in  $(k-I)N^{L-2}$  guesses
      - ► Etc.



## Use specific cracking algorithm to compute number of guesses it would take



## Guessing Passwords

- Approach one: Markov Chain
  - For each character the probability of the next character varies
  - First guess highest probability first char
  - Next guess highest probability subsequent character
  - Repeat
  - If fail, go to next highest probability character and continue



## Guessing Passwords

- Suppose password is "CAC"
  - In character set {ABC}
- Start with highest probability start A
  - Compute all passwords that start with A
  - In highest probability order count so far  $k^n = 9$
- Then go to the next highest prob. start say C
  - Next highest prob. for second char A
  - Then A, B, C for third char
- For a guess number of 11





## Guess Number

- to crack a specific password
  - Produce a deterministic guess ordering
- For "Weir" cracker
  - (Probabilistic Context-Free Grammar)
  - Uses probabilities of password structures
    - E.g., Small letter ^ N + Number ^ I + Capital letter ^ M ...
- Computing guess number
  - Determine the guesses necessary to reach the "probability group" for that password
  - Add number of further guesses to reach exact password



## Use specific cracking algorithm to compute number of guesses it would take







## Guessing Passwords

- Suppose highest password is "BAI"
  - In character set {ABI}
- Start with highest probability struct {L<sup>2</sup>D<sup>1</sup>}
  - Search for most likely L<sup>2</sup> and most likely D<sup>1</sup>
- For Markov, search from highest probability A
  - $K^n = 2$
  - Next highest prob. B
  - Then A
  - Then | for D<sup>1</sup>
- For a guess number of 3



## How Many Guesses For?

## By password-composition policy



most comprehensive training set, which combines our passwords with public data.

![](_page_22_Picture_6.jpeg)

Figure 1. The number of passwords cracked vs. number of guesses, per condition, for experiment E. This experiment uses the Weir calculator and our

![](_page_22_Picture_9.jpeg)

## Something you have ...

- Tokens (transponders, ...)
  - Speedpass, EZ-pass
  - SecureID
- Smartcards
  - Unpowered processors
  - Small NV storage
  - Tamper resistant

- Digital Certificates (used by Websites to authenticate themselves to customers)
  - More on this later ...

![](_page_23_Picture_11.jpeg)

![](_page_23_Picture_12.jpeg)

![](_page_23_Picture_13.jpeg)

## A (simplified) sample token device

- A one-time password system that essentially uses a hash chain as authenticators.
  - For seed (S) and chain length (I)
  - Tamperproof token encodes S in firmware

- Device display shows password for epoch i
- Time synchronization allows authentication server to know what i is expected, and authenticate the user.
- Note: somebody can see your token display at some time but learn nothing useful for later periods.

![](_page_24_Picture_9.jpeg)

![](_page_24_Picture_10.jpeg)

![](_page_24_Picture_13.jpeg)

![](_page_24_Picture_14.jpeg)

![](_page_24_Figure_16.jpeg)

![](_page_24_Picture_17.jpeg)

![](_page_24_Picture_18.jpeg)

## Something your are ...

- Biometrics measure some physical characteristic
  - Fingerprint, face recognition, retina scanners, voice, signature, DNA
  - Can be extremely accurate and fast
  - Active biometrics authenticate
  - Passive biometrics recognize

- Issues with biometrics?
  - Revocation lost fingerprint?
  - "fuzzy" credential, e.g., your face changes based on mood ...
  - Great for physical security, not feasible for on-line systems

![](_page_25_Picture_11.jpeg)

![](_page_25_Picture_14.jpeg)

![](_page_25_Picture_16.jpeg)

![](_page_25_Picture_17.jpeg)

![](_page_25_Picture_18.jpeg)

## Biometrics Example

- A fingerprint biometric device (of several)

![](_page_26_Picture_4.jpeg)

![](_page_26_Picture_6.jpeg)

## record the conductivity of the surface of your finger to build a "map" of the ridges scanned map converted into a graph by looking for landmarks, e.g., ridges, cores, ...

![](_page_26_Figure_9.jpeg)

![](_page_26_Picture_10.jpeg)

A fingerprint is represented in the form of a graph whose nodes correspond to ridges in the print. Edges of the graph connect nodes that represent neighboring or

## Fingerprint Biometrics (cont.)

- Graph is compared to database of authentic identities
- Graph is same, the person deemed "authentic"
  - This is a variant of the graph isomorphism problem
  - Problem: what does it mean to be the "same enough"
    - rotation
    - imperfect contact
    - finger damage

## • Fundamental Problem: False accept vs. false reject rates?

![](_page_27_Picture_10.jpeg)

![](_page_27_Picture_11.jpeg)